

# EECS498-003

# Formal Verification of Systems Software

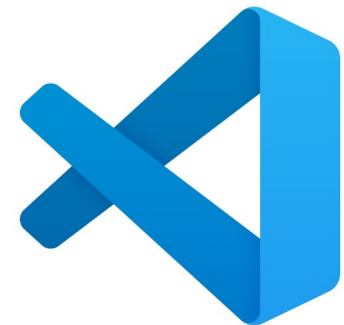
Material and slides created by  
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PREVIOUSLY ON  
FORMAL VERIFICATION

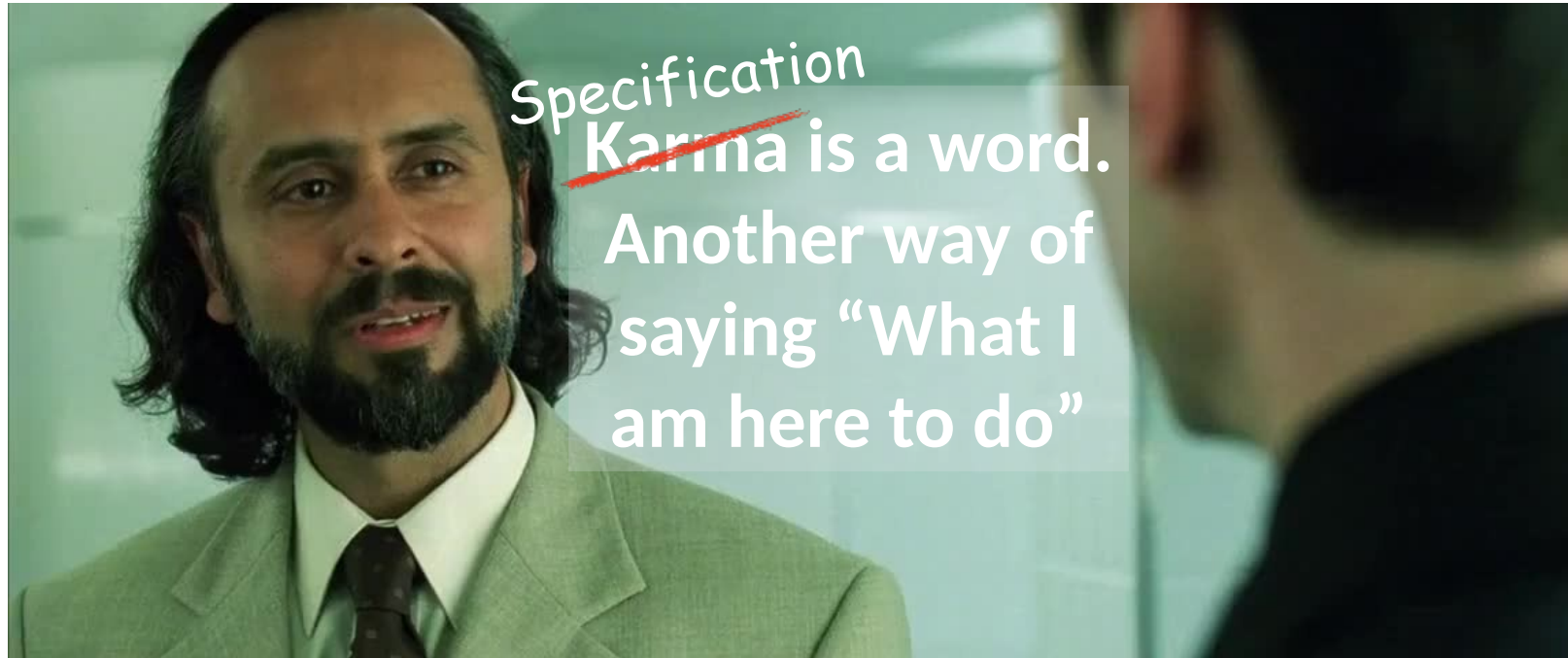
# Recursion: exporting ensures

```
function Evens(count:int) : (outseq:seq<int>)  
  ensures forall idx :: 0<=idx<|outseq| ==> outseq[idx] == 2 * idx  
{  
  if count==0 then [] else Evens(count) + [2 * (count-1)]  
}
```



VSCode transition

# Chapter 2: Specification



# How to specify our programs

Attempt #1: Just tell your programmers what you want them to code

*Writing is nature's way of letting you know how sloppy your thinking is*  
-Dick Guindon

# How to specify our programs

Attempt #2: Write down an English description (aka a design doc)

*Mathematics is nature's way of letting you know how sloppy your writing is*

-Leslie Lamport

*Formal mathematics is nature's way of letting you know how sloppy your mathematics is*

-Leslie Lamport

# Formal specification

A way to define formally (i.e. precisely) what your program should do

Before you start writing code, make sure you know what code is supposed to be doing

Before you start writing a proof, make sure you know what you are proving

# Specification

A specification defines *which executions are allowable*

```
lemma Double(x:int) returns (y:int)
  ensures y == 2*x
{
  ...
}
```

(x=1, y=2)	✓
(x=2, y=4)	✓
(x=2, y=2)	✗
(x=-3, y=-6)	✓
(x=-2, y=4)	✗



# Ways to specify what the program should do

- C-style assertions

```
y = 2*x;  
assert(y==2*x)
```

- Postconditions

```
lemma Double(x:int) returns  
(y:int)  
  ensures y == 2*x  
{  
  y := 2*x;  
}
```

- Properties/invariants

“At most one node holds the lock at any time”

- Refinement

- Linearizability
- Equivalence to logically centralized service

# Specification is trusted

Formal verification: proving that your protocol or implementation meets the spec

*You cannot prove that the spec is correct*

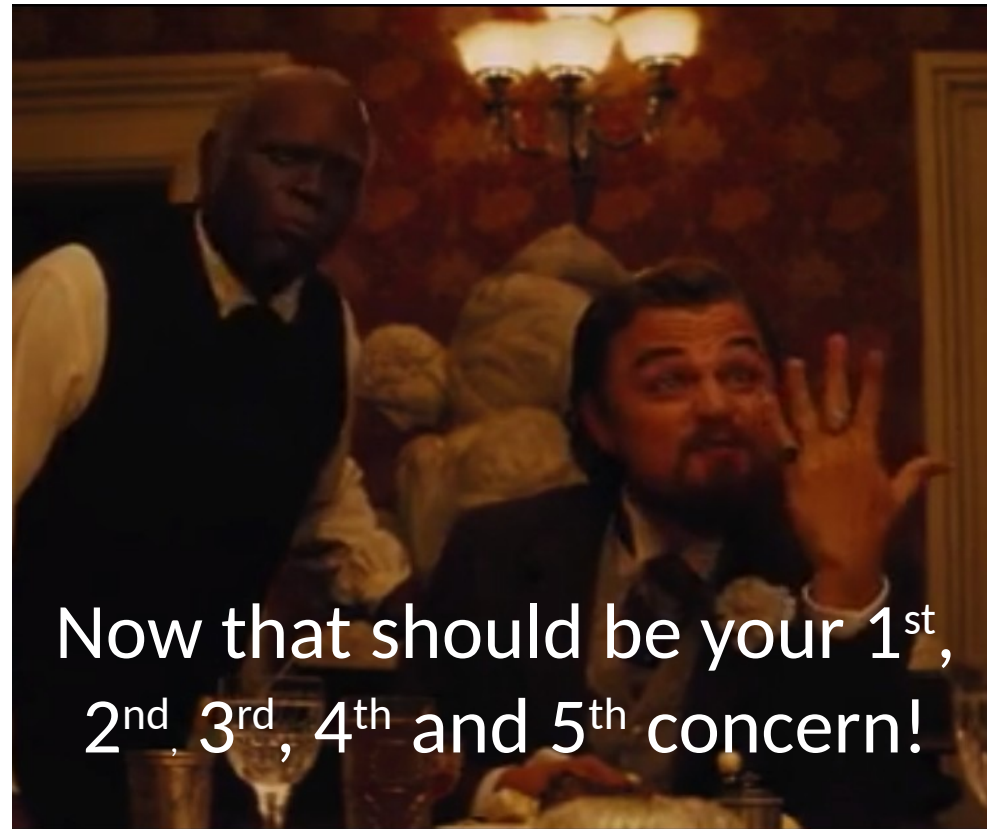
You have to *trust* your spec

Your proof is as good as your spec

A wrong spec is one of the few ways to introduce bugs into formally verified code

# Check your spec

1. Check your spec!
2. Check your spec!
3. Check your spec!
4. Check your spec!
5. Check your spec!



# Research detour: IronSpec (OSDI '24)

- A methodology and tool to find bugs in formal specifications
- Found multiple spec bugs in real-world Dafny codebases

## **IronSpec: Increasing the Reliability of Formal Specifications**

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# The benefit of specification

The spec is typically *much smaller* than the code

- So we have to inspect a few lines of code only

Dijkstra's algorithm spec

```
IsShortestPath(g, p) {  
    && IsPath(g, p)  
    && forall p2 :: IsPath(g, p2) ==> |p| <= |p2|  
}
```

# A good spec

A good spec is *correct/complete*

- It precludes all undesirable behaviors

Example: IsMaxIndex

```
predicate IsMaxIndex(a:seq<int>, x:int) {  
  && 0 < x < |a|  
  && (forall i | 0 < i < |a| :: a[i] <= a[x])  
}
```

# A good spec (cont.)

A good spec is *concise*

- It elides every irrelevant concept
- Is simple and easy to read

```
predicate IsMaxIndex(a:seq<int>, x:int) {  
  && 0 <= x < |a|  
  && (forall i | 0 <= i < |a| :: a[i] <= a[x])  
}
```

# A good spec (cont.)

A good spec is *abstract*

- It doesn't constrain the implementation

## Dijkstra's algorithm spec

```
IsShortestPath(g:Graph, p:Path) {  
    && IsPath(g, p)  
    && forall p2 :: IsPath(g, p2) ==> |p| <= |p2|  
}
```



# Edsger W. Dijkstra



- 1972 Turing Award winner
- Inventor of:
  - Dijkstra's shortest path algorithm
  - Semaphores
  - The THE operating system
  - Banker's algorithm
- *“Progress is possible only if we train ourselves to think about programs without thinking of them as pieces of executable code.”*

# Verification and the “eradication” of bugs

Frequent quote from verification experts

- “We prove that there are no bugs at all...”

Frequent quote from verification skeptics

- “Nonsense! You can still have bugs in your spec”

The truth is somewhere in the middle

- Yes, your spec may have bugs
- But do you prefer inspecting 30 lines for bugs or 30000?

# Autograder submissions

- Reminder: PS1 deadline is September 19
- Dafny timeouts/out of resource
  - `dafny /trace`
  - `dafny /rlimit:32700000`
- Remember to pull the code repo for the starter code of examples presented in class

# Some new Dafny syntax

## Datatype member functions

```
datatype Pet = Dog | Cat | Ant | Spider {  
  function CountLegs() : int {  
    match this  
      case Dog => 4  
      case Cat => 4  
      case Ant => 6  
      case Spider => 8  
    }  
  }  
}  
  
function ShoesForTwo(pet: Pet) : int {  
  2 * pet.CountLegs()  
}
```

# Some new Dafny syntax

## Calc statements

```
assert a == b;  
assert b == c;  
assert c == d;
```

```
calc {  
  a;  
  b;  
  c;  
  d;  
}
```

# Some new Dafny syntax

## Calc statements

```
assert a == b;  
assert b == c;  
assert c == d;
```

```
calc {  
  a;  
  { MyUsefulLemma(a,b); }  
  b;  
  c;  
  d;  
}
```

# Some new Dafny syntax

## Calc statements

```
assert a == b;  
assert b == c;  
assert c == d;
```

```
calc ==> {  
  a;  
  { MyUsefulLemma(a,b); }  
  b;  
  c;  
  d;  
}
```

# Some new Dafny syntax

Choose operator

```
assert 1 % 7 == 1;  
assert exists x :: x % 7 == 1;  
var x :| x % 7 == 1;
```

Choose x such that...